

Commutative Algebra and Approximation Theory

Michael DiPasquale
Oklahoma State University
Colloquium

Piecewise Polynomials

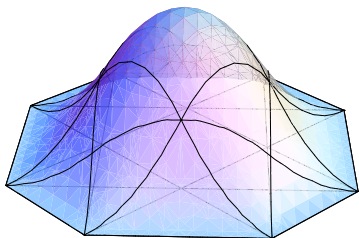
Spline

A piecewise polynomial function, continuously differentiable to some order.

Piecewise Polynomials

Spline

A piecewise polynomial function, continuously differentiable to some order.



The Zwart-Powell element, a C^1 spline of degree 2

Univariate Splines

Most widely studied case: approximation of a function $f(x)$ over an interval $\Delta = [a, b] \subset \mathbb{R}$ by C^r piecewise polynomials.

Most widely studied case: approximation of a function $f(x)$ over an interval $\Delta = [a, b] \subset \mathbb{R}$ by C^r piecewise polynomials.

- Subdivide $\Delta = [a, b]$ into subintervals:
$$\Delta = [a_0, a_1] \cup [a_1, a_2] \cup \cdots \cup [a_{n-1}, a_n]$$
- Find a basis for the vector space $C_d^r(\Delta)$ of C^r piecewise polynomial functions on Δ with degree at most d (B-splines!)
- Find best approximation to $f(x)$ in $C_d^r(\Delta)$

Two Subintervals

$\Delta = [a_0, a_1] \cup [a_1, a_2]$ (assume WLOG $a_1 = 0$)

$$(f_1, f_2) \in C_d^r(\Delta) \iff f_1^{(i)}(0) = f_2^{(i)}(0) \text{ for } 0 \leq i \leq r$$

$$\iff x^{r+1} | (f_2 - f_1)$$

$$\iff (f_2 - f_1) \in \langle x^{r+1} \rangle$$

Two Subintervals

$\Delta = [a_0, a_1] \cup [a_1, a_2]$ (assume WLOG $a_1 = 0$)

$$\begin{aligned}(f_1, f_2) \in C_d^r(\Delta) &\iff f_1^{(i)}(0) = f_2^{(i)}(0) \text{ for } 0 \leq i \leq r \\ &\iff x^{r+1} | (f_2 - f_1) \\ &\iff (f_2 - f_1) \in \langle x^{r+1} \rangle\end{aligned}$$

Even more explicitly:

- $f_1(x) = b_0 + b_1x + \cdots + b_dx^d$
- $f_2(x) = c_0 + c_1x + \cdots + c_dx^d$
- $(f_0, f_1) \in C_d^r(\Delta) \iff b_0 = c_0, \dots, b_r = c_r.$

Two Subintervals

$\Delta = [a_0, a_1] \cup [a_1, a_2]$ (assume WLOG $a_1 = 0$)

$$\begin{aligned}(f_1, f_2) \in C_d^r(\Delta) &\iff f_1^{(i)}(0) = f_2^{(i)}(0) \text{ for } 0 \leq i \leq r \\ &\iff x^{r+1} | (f_2 - f_1) \\ &\iff (f_2 - f_1) \in \langle x^{r+1} \rangle\end{aligned}$$

Even more explicitly:

- $f_1(x) = b_0 + b_1x + \cdots + b_dx^d$
- $f_2(x) = c_0 + c_1x + \cdots + c_dx^d$
- $(f_0, f_1) \in C_d^r(\Delta) \iff b_0 = c_0, \dots, b_r = c_r.$

$$\dim C_d^r(\Delta) = \begin{cases} d + 1 & \text{if } d \leq r \\ (d + 1) + (d - r) & \text{if } d > r \end{cases}$$

Note: $\dim C_d^r(\Delta)$ is polynomial in d for $d > r$.

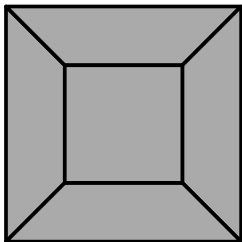
More General Problem: Compute $\dim C_d^r(\Delta)$ where $\Delta \subset \mathbb{R}^n$ is

More General Problem: Compute $\dim C_d^r(\Delta)$ where $\Delta \subset \mathbb{R}^n$ is

- a **polytopal complex**
- pure n -dimensional
- a **pseudomanifold**

More General Problem: Compute $\dim C_d^r(\Delta)$ where $\Delta \subset \mathbb{R}^n$ is

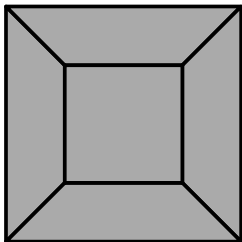
- a **polytopal complex**
- pure n -dimensional
- a **pseudomanifold**



A polytopal complex

More General Problem: Compute $\dim C_d^r(\Delta)$ where $\Delta \subset \mathbb{R}^n$ is

- a **polytopal complex**
- pure n -dimensional
- a **pseudomanifold**



A polytopal complex

(Algebraic) Spline Criterion:

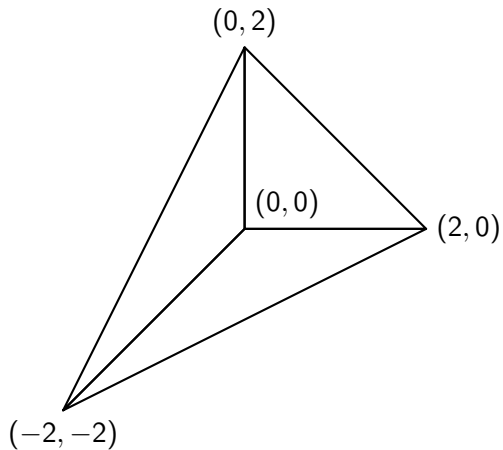
- If $\tau \in \Delta_{n-1}$, $l_\tau =$ affine form vanishing on affine span of τ
- Collection $\{f_\sigma\}_{\sigma \in \Delta_n}$ glue to $F \in C^r(\Delta) \iff$ for every pair of adjacent facets $\sigma_1, \sigma_2 \in \Delta_n$ with $\sigma_1 \cap \sigma_2 = \tau \in \Delta_{n-1}$, $l_\tau^{r+1} \mid (f_{\sigma_1} - f_{\sigma_2})$

Who Cares?

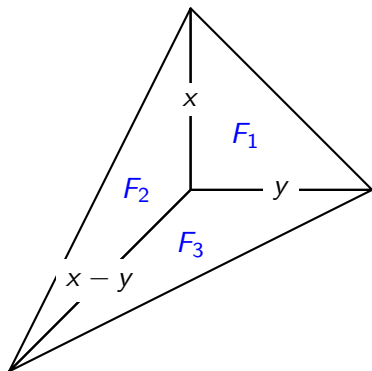
- 1 Computation of $\dim C_d^r(\Delta)$ for higher dimensions initiated by [Strang '73] in connection with finite element method
- 2 Data fitting in approximation theory
- 3 [Farin '97] Computer Aided Geometric Design (CAGD) - building surfaces by splines.
- 4 [Payne '06] Toric Geometry - Equivariant cohomology rings of toric varieties are rings of continuous splines on the fan (under appropriate conditions).

Part I: Continuous Splines and Freeness

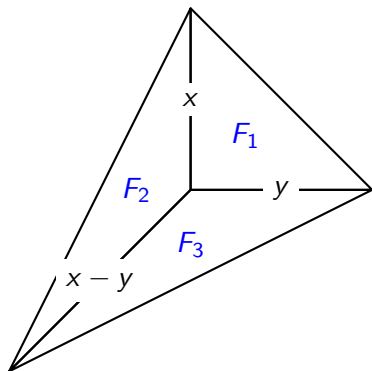
Continuous Splines



Continuous Splines



Continuous Splines

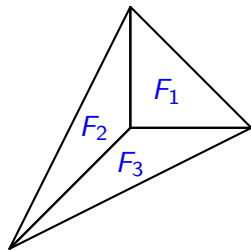


$$(F_1, F_2, F_3) \in C^0(\Delta) \iff \\ \exists f_1, f_2, f_3 \text{ so that}$$

$$F_1 - F_2 = f_1 x$$

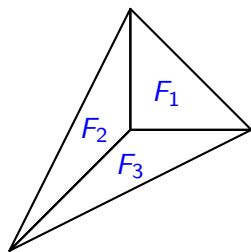
$$F_2 - F_3 = f_2(x - y)$$

$$F_3 - F_1 = f_3 y$$



$(F_1, F_2, F_3) \in C^0(\Delta) \iff$ there are
 f_1, f_2, f_3 so that

Spline Matrix



$(F_1, F_2, F_3) \in C^0(\Delta) \iff$ there are f_1, f_2, f_3 so that

$$\begin{pmatrix} 1 & -1 & 0 & x & 0 & 0 \\ 0 & 1 & -1 & 0 & x-y & 0 \\ -1 & 0 & 1 & 0 & 0 & y \end{pmatrix} \begin{pmatrix} F_1 \\ F_2 \\ F_3 \\ -f_1 \\ -f_2 \\ -f_3 \end{pmatrix} = 0$$

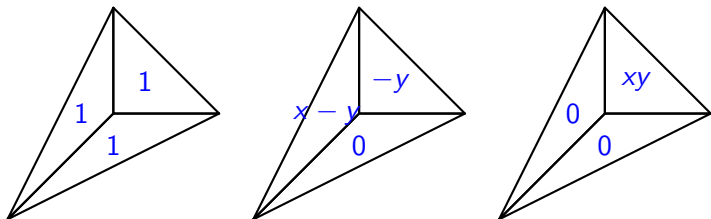
This matrix constructed in [Billera-Rose '91].

Observations

- $C^0(\Delta)$, the kernel of this matrix, is a **graded** $\mathbb{R}[x, y]$ -module (matrix entries are homogeneous).
- $C^0(\Delta)_d :=$ splines of degree d

Observations

- $C^0(\Delta)$, the kernel of this matrix, is a **graded** $\mathbb{R}[x, y]$ -module (matrix entries are homogeneous).
- $C^0(\Delta)_d :=$ splines of degree d
- Every spline in $C^0(\Delta)$ can be written uniquely as a polynomial combination of the three splines pictured below:



Observations, continued

$C^0(\Delta)$ is a **free** $R = \mathbb{R}[x, y]$ -module generated in degrees 0,1,2.
Record degrees as $C^0(\Delta) \cong R \oplus R(-1) \oplus R(-2)$.

Observations, continued

$C^0(\Delta)$ is a **free** $R = \mathbb{R}[x, y]$ -module generated in degrees 0,1,2.
Record degrees as $C^0(\Delta) \cong R \oplus R(-1) \oplus R(-2)$.

$$\dim C^0(\Delta)_d = \binom{d+1}{1} + \binom{(d+1)-1}{1} + \binom{(d+1)-2}{1}$$

Observations, continued

$C^0(\Delta)$ is a **free** $R = \mathbb{R}[x, y]$ -module generated in degrees 0,1,2.
Record degrees as $C^0(\Delta) \cong R \oplus R(-1) \oplus R(-2)$.

$$\begin{aligned} \dim C^0(\Delta)_d &= \binom{d+1}{1} + \binom{(d+1)-1}{1} + \binom{(d+1)-2}{1} \\ &= 3d \text{ for } d \geq 1 \end{aligned}$$

Observations, continued

$C^0(\Delta)$ is a **free** $R = \mathbb{R}[x, y]$ -module generated in degrees 0,1,2.
Record degrees as $C^0(\Delta) \cong R \oplus R(-1) \oplus R(-2)$.

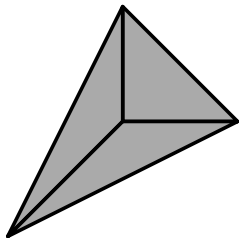
$$\begin{aligned}\dim C^0(\Delta)_d &= \binom{d+1}{1} + \binom{(d+1)-1}{1} + \binom{(d+1)-2}{1} \\ &= 3d \text{ for } d \geq 1\end{aligned}$$

$$\begin{aligned}\dim C^0(\Delta)_d &= \\ \dim C^0(\widehat{\Delta})_d &= \binom{d+2}{2} + \binom{(d+2)-1}{2} + \binom{(d+2)-2}{2} \\ &= \frac{3}{2}d^2 + \frac{3}{2}d + 1 \text{ for } d \geq 0,\end{aligned}$$

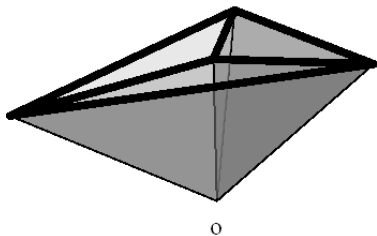
where $\widehat{\Delta}$ is the cone over Δ .

Coning Construction

- $\widehat{\Delta} \subset \mathbb{R}^{n+1}$ denotes the cone over $\Delta \subset \mathbb{R}^n$.



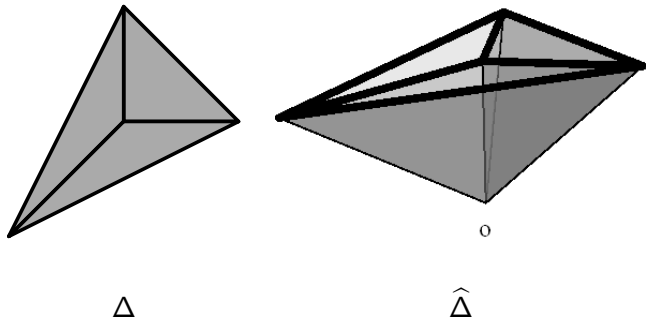
Δ



$\widehat{\Delta}$

Coning Construction

- $\widehat{\Delta} \subset \mathbb{R}^{n+1}$ denotes the cone over $\Delta \subset \mathbb{R}^n$.



- $C^r(\widehat{\Delta})$ is always a **graded** algebra over $S = \mathbb{R}[x_0, \dots, x_n]$
- $C_d^r(\Delta) \cong C^r(\widehat{\Delta})_d$ [Billera-Rose '91]

- Freeness of $C^r(\widehat{\Delta}) \implies$ straightforward computation of $\dim C_d^r(\Delta)$.

Consequences of Freeness

- Freeness of $C^r(\widehat{\Delta}) \implies$ straightforward computation of $\dim C_d^r(\Delta)$.
- [Schenck-Stillman '97] Many widely-used partitions Δ actually satisfy the property that $C^r(\widehat{\Delta})$ is free (type I and II triangulations, cross-cut partitions, rectangular meshes, etc.)

Consequences of Freeness

- Freeness of $C^r(\widehat{\Delta}) \implies$ straightforward computation of $\dim C_d^r(\Delta)$.
- [Schenck-Stillman '97] Many widely-used partitions Δ actually satisfy the property that $C^r(\widehat{\Delta})$ is free (type I and II triangulations, cross-cut partitions, rectangular meshes, etc.)
- [Billera-Rose '92] criteria for freeness in terms of localization
- [Yuzvinsky '92] criteria for freeness in terms of sheaves on posets
- [Schenck '97] criteria for freeness in terms of homologies of a chain complex (Δ simplicial)

Face Rings of Simplicial Complexes

Face Ring of Δ

Δ a simplicial complex.

$$A_{\Delta} = \mathbb{R}[x_v \mid v \text{ a vertex of } \Delta] / I_{\Delta},$$

where I_{Δ} is the ideal generated by monomials corresponding to non-faces.

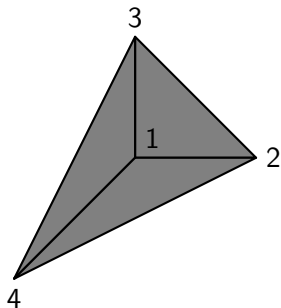
Face Rings of Simplicial Complexes

Face Ring of Δ

Δ a simplicial complex.

$$A_{\Delta} = \mathbb{R}[x_v \mid v \text{ a vertex of } \Delta] / I_{\Delta},$$

where I_{Δ} is the ideal generated by monomials corresponding to non-faces.



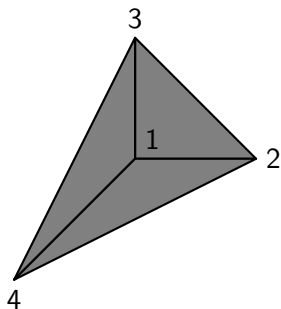
Face Rings of Simplicial Complexes

Face Ring of Δ

Δ a simplicial complex.

$$A_{\Delta} = \mathbb{R}[x_v \mid v \text{ a vertex of } \Delta] / I_{\Delta},$$

where I_{Δ} is the ideal generated by monomials corresponding to non-faces.



- Nonfaces are $\{1, 2, 3, 4\}, \{2, 3, 4\}$
- $I_{\Delta} = \langle x_2 x_3 x_4 \rangle$
- $A_{\Delta} = \mathbb{R}[x_1, x_2, x_3, x_4] / I_{\Delta}$

C^0 for Simplicial Splines [Billera '89]

- $C^0(\widehat{\Delta}) \cong A_{\Delta}$, the face ring of Δ .

C^0 for Simplicial Splines [Billera '89]

- $C^0(\widehat{\Delta}) \cong A_{\Delta}$, the face ring of Δ .
- $\dim C_d^0(\Delta) = \sum_{i=0}^n f_i \binom{d-1}{i}$ for $d > 0$, where $f_i = \#i$ -faces of Δ .

C^0 for Simplicial Splines [Billera '89]

- $C^0(\widehat{\Delta}) \cong A_{\Delta}$, the face ring of Δ .
- $\dim C_d^0(\Delta) = \sum_{i=0}^n f_i \binom{d-1}{i}$ for $d > 0$, where $f_i = \#i$ -faces of Δ .

Moreover, if Δ is homeomorphic to a disk, then $C^0(\widehat{\Delta})$ is free.

Nonsimplicial Case

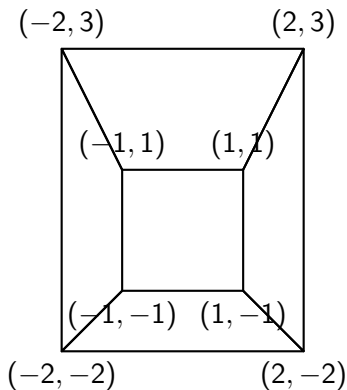
Nonfreeness for Polytopal Complexes [D. '12]

$C^0(\widehat{\Delta})$ need not be free if Δ has nonsimplicial faces.

Nonsimplicial Case

Nonfreeness for Polytopal Complexes [D. '12]

$C^0(\widehat{\Delta})$ need not be free if Δ has nonsimplicial faces.



$C^0(\widehat{\Delta})$ is **not** a free module over $\mathbb{R}[x, y, z]$.

Part II: Hilbert Polynomials and Regularity

Some Graded Commutative Algebra

Given a finitely generated graded $S = \mathbb{R}[x_1, \dots, x_n]$ -module M (like $C^r(\widehat{\Delta})$).

- $HF(M, d) := \dim M_d$ is the **Hilbert function** of M .

Some Graded Commutative Algebra

Given a finitely generated graded $S = \mathbb{R}[x_1, \dots, x_n]$ -module M (like $C^r(\widehat{\Delta})$).

- $HF(M, d) := \dim M_d$ is the **Hilbert function** of M .
- If $d \gg 0$, $HF(M, d) = HP(M, d)$, where $HP(M, d)$ is the **Hilbert polynomial** of M .

Some Graded Commutative Algebra

Given a finitely generated graded $S = \mathbb{R}[x_1, \dots, x_n]$ -module M (like $C^r(\widehat{\Delta})$).

- $HF(M, d) := \dim M_d$ is the **Hilbert function** of M .
- If $d \gg 0$, $HF(M, d) = HP(M, d)$, where $HP(M, d)$ is the **Hilbert polynomial** of M .
- Upshot: $\dim C_d^r(\Delta) = \dim C^r(\widehat{\Delta})_d$ is eventually polynomial in d (in fact, linear combination of binomial coefficients)

The Good News and the Bad News

Good news: $HP(C^r(\widehat{\Delta}), d)$ has been computed for $\Delta \subset \mathbb{R}^2$.

- Δ simplicial: [Alfeld-Schumaker '90, Hong '91], [Ibrahim-Schumaker '91]
- Δ nonsimplicial: [McDonald-Schenck '09]

The Good News and the Bad News

Good news: $HP(C^r(\widehat{\Delta}), d)$ has been computed for $\Delta \subset \mathbb{R}^2$.

- Δ simplicial: [Alfeld-Schumaker '90, Hong '91], [Ibrahim-Schumaker '91]
- Δ nonsimplicial: [McDonald-Schenck '09]

Bad news: $\dim C_d^r(\Delta)$ is still a mystery for small d .

- $\dim C_3^1(\Delta)$ still unknown for $\Delta \subset \mathbb{R}^2$!

Planar Hilbert Polynomials

Planar Simplicial Dimension [Alfeld-Schumaker '90]

If $\Delta \subset \mathbb{R}^2$ is a simply connected triangulation and $d \geq 3r + 1$, then

$$\dim C_d^r(\Delta) = \binom{d+2}{2} + \binom{d-r+1}{2} f_1^0 - \left(\binom{d+2}{2} - \binom{r+2}{2} \right) f_0^0 + \sigma,$$

Planar Hilbert Polynomials

Planar Simplicial Dimension [Alfeld-Schumaker '90]

If $\Delta \subset \mathbb{R}^2$ is a simply connected triangulation and $d \geq 3r + 1$, then

$$\dim C_d^r(\Delta) = \binom{d+2}{2} + \binom{d-r+1}{2} f_1^0 - \left(\binom{d+2}{2} - \binom{r+2}{2} \right) f_0^0 + \sigma,$$

- f_i^0 is the number of interior i -dimensional faces.

Planar Hilbert Polynomials

Planar Simplicial Dimension [Alfeld-Schumaker '90]

If $\Delta \subset \mathbb{R}^2$ is a simply connected triangulation and $d \geq 3r + 1$, then

$$\dim C_d^r(\Delta) = \binom{d+2}{2} + \binom{d-r+1}{2} f_1^0 - \left(\binom{d+2}{2} - \binom{r+2}{2} \right) f_0^0 + \sigma,$$

- f_i^0 is the number of interior i -dimensional faces.
- $n(v_i) = \#$ distinct slopes at an interior vertex v_i .

Planar Hilbert Polynomials

Planar Simplicial Dimension [Alfeld-Schumaker '90]

If $\Delta \subset \mathbb{R}^2$ is a simply connected triangulation and $d \geq 3r + 1$, then

$$\dim C_d^r(\Delta) = \binom{d+2}{2} + \binom{d-r+1}{2} f_1^0 - \left(\binom{d+2}{2} - \binom{r+2}{2} \right) f_0^0 + \sigma,$$

- f_i^0 is the number of interior i -dimensional faces.
- $n(v_i) = \#$ distinct slopes at an interior vertex v_i .
- $\sigma_i = \sum_j \max\{(r+1+j(1-n(v_i))), 0\}$.

Planar Hilbert Polynomials

Planar Simplicial Dimension [Alfeld-Schumaker '90]

If $\Delta \subset \mathbb{R}^2$ is a simply connected triangulation and $d \geq 3r + 1$, then

$$\dim C_d^r(\Delta) = \binom{d+2}{2} + \binom{d-r+1}{2} f_1^0 - \left(\binom{d+2}{2} - \binom{r+2}{2} \right) f_0^0 + \sigma,$$

- f_i^0 is the number of interior i -dimensional faces.
- $n(v_i) = \#$ distinct slopes at an interior vertex v_i .
- $\sigma_i = \sum_j \max\{(r+1+j(1-n(v_i))), 0\}$.
- $\sigma = \sum \sigma_i$.

Planar Hilbert Polynomials

Planar Simplicial Dimension [Alfeld-Schumaker '90]

If $\Delta \subset \mathbb{R}^2$ is a simply connected triangulation and $d \geq 3r + 1$, then

$$\dim C_d^r(\Delta) = \binom{d+2}{2} + \binom{d-r+1}{2} f_1^0 - \left(\binom{d+2}{2} - \binom{r+2}{2} \right) f_0^0 + \sigma,$$

- f_i^0 is the number of interior i -dimensional faces.
- $n(v_i) = \#$ distinct slopes at an interior vertex v_i .
- $\sigma_i = \sum_j \max\{(r+1+j(1-n(v_i))), 0\}$.
- $\sigma = \sum \sigma_i$.

Conjecture [Schenck '97]

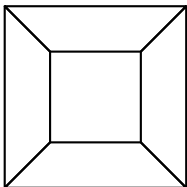
Above formula holds for $d \geq 2r + 1$.

Planar Hilbert Polynomials

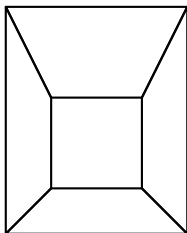
- $\Delta \subset \mathbb{R}^2$ a simply connected polytopal complex
- [McDonald-Schenck '09] give formulas for coefficients of $HP(C^r(\widehat{\Delta}), d)$

Planar Hilbert Polynomials

- $\Delta \subset \mathbb{R}^2$ a simply connected polytopal complex
- [McDonald-Schenck '09] give formulas for coefficients of $HP(C^r(\widehat{\Delta}), d)$



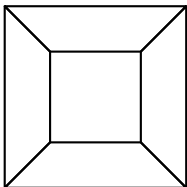
$$HP(C^0(\widehat{\Delta}), d) = \frac{5}{2}d^2 - \frac{1}{2}d + 2$$



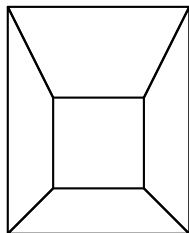
$$HP(C^0(\widehat{\Delta}), d) = \frac{5}{2}d^2 - \frac{1}{2}d + 1$$

Planar Hilbert Polynomials

- $\Delta \subset \mathbb{R}^2$ a simply connected polytopal complex
- [McDonald-Schenck '09] give formulas for coefficients of $HP(C^r(\widehat{\Delta}), d)$



$$HP(C^0(\widehat{\Delta}), d) = \frac{5}{2}d^2 - \frac{1}{2}d + 2$$

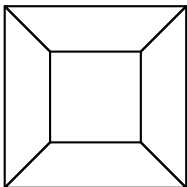


$$HP(C^0(\widehat{\Delta}), d) = \frac{5}{2}d^2 - \frac{1}{2}d + 1$$

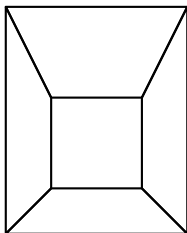
How large does d have to be for $\dim C_d^r(\Delta) = HP(C^r(\widehat{\Delta}), d)$?

Planar Hilbert Polynomials

- $\Delta \subset \mathbb{R}^2$ a simply connected polytopal complex
- [McDonald-Schenck '09] give formulas for coefficients of $HP(C^r(\hat{\Delta}), d)$



$$HP(C^0(\hat{\Delta}), d) = \frac{5}{2}d^2 - \frac{1}{2}d + 2$$



$$HP(C^0(\hat{\Delta}), d) = \frac{5}{2}d^2 - \frac{1}{2}d + 1$$

How large does d have to be for $\dim C_d^r(\Delta) = HP(C^r(\hat{\Delta}), d)$?

In simplicial case, $d \geq 3r + 1$ suffices.

Large degree generators

Proposition [D. '14]

Given an n -polytope $A \subset \mathbb{R}^n$ and a choice of codimension 1 face $\tau \in A_{n-1}$, there is a polytopal complex $\mathcal{P}(A)$ having A as a facet so that

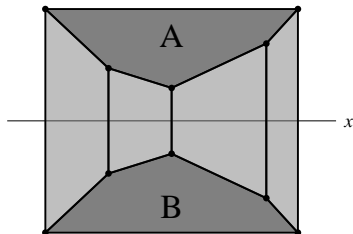
- 1 Every codimension 1 face of A except τ is interior to $\mathcal{P}(A)$
- 2 There is a minimal generator of $C^r(\widehat{\mathcal{P}(A)})$ supported only on A .

Large degree generators

Proposition [D. '14]

Given an n -polytope $A \subset \mathbb{R}^n$ and a choice of codimension 1 face $\tau \in A_{n-1}$, there is a polytopal complex $\mathcal{P}(A)$ having A as a facet so that

- 1 Every codimension 1 face of A except τ is interior to $\mathcal{P}(A)$
- 2 There is a minimal generator of $C^r(\widehat{\mathcal{P}(A)})$ supported only on A .



$C^r(\widehat{\Delta})$ has minimal generator of degree $4(r + 1)$

A Positive Result

Agreement of Hilbert Function and Polynomial [D. '14]

$\Delta \subset \mathbb{R}^2$ a planar polytopal complex. Let $F =$ maximum number of edges of a polygon of Δ . Then

$$HP(C^r(\widehat{\Delta}), d) = \dim C_d^r(\Delta) \text{ for } d \geq (2F - 1)(r + 1) - 1$$

A Positive Result

Agreement of Hilbert Function and Polynomial [D. '14]

$\Delta \subset \mathbb{R}^2$ a planar polytopal complex. Let $F =$ maximum number of edges of a polygon of Δ . Then

$$HP(C^r(\widehat{\Delta}), d) = \dim C_d^r(\Delta) \text{ for } d \geq (2F - 1)(r + 1) - 1$$

This is the first such result for nonsimplicial complexes.

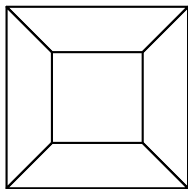
A Positive Result

Agreement of Hilbert Function and Polynomial [D. '14]

$\Delta \subset \mathbb{R}^2$ a planar polytopal complex. Let $F =$ maximum number of edges of a polygon of Δ . Then

$$HP(C^r(\widehat{\Delta}), d) = \dim C_d^r(\Delta) \text{ for } d \geq (2F - 1)(r + 1) - 1$$

This is the first such result for nonsimplicial complexes.



$$\begin{aligned} & HP(C^0(\widehat{\Delta}), d) \\ &= \frac{5}{2}d^2 - \frac{1}{2}d + 2 \end{aligned}$$

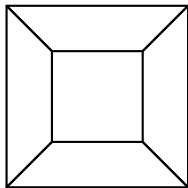
A Positive Result

Agreement of Hilbert Function and Polynomial [D. '14]

$\Delta \subset \mathbb{R}^2$ a planar polytopal complex. Let $F =$ maximum number of edges of a polygon of Δ . Then

$$HP(C^r(\widehat{\Delta}), d) = \dim C_d^r(\Delta) \text{ for } d \geq (2F - 1)(r + 1) - 1$$

This is the first such result for nonsimplicial complexes.



- $F = 4$
- $\implies \dim C_d^0(\Delta) = \frac{5}{2}d^2 - \frac{1}{2}d + 2$ for $d \geq 6$

$$\begin{aligned} & HP(C^0(\widehat{\Delta}), d) \\ &= \frac{5}{2}d^2 - \frac{1}{2}d + 2 \end{aligned}$$

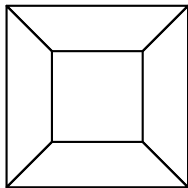
A Positive Result

Agreement of Hilbert Function and Polynomial [D. '14]

$\Delta \subset \mathbb{R}^2$ a planar polytopal complex. Let $F =$ maximum number of edges of a polygon of Δ . Then

$$HP(C^r(\widehat{\Delta}), d) = \dim C_d^r(\Delta) \text{ for } d \geq (2F - 1)(r + 1) - 1$$

This is the first such result for nonsimplicial complexes.



$$\begin{aligned} &HP(C^0(\widehat{\Delta}), d) \\ &= \frac{5}{2}d^2 - \frac{1}{2}d + 2 \end{aligned}$$

- $F = 4$
- $\implies \dim C_d^0(\Delta) = \frac{5}{2}d^2 - \frac{1}{2}d + 2$ for $d \geq 6$
- Macaulay2:
 $\dim C_d^0(\Delta) = \frac{5}{2}d^2 - \frac{1}{2}d + 2$ for $d \geq 1$

The Technique: Regularity

Set $S = \mathbb{R}[x_1, \dots, x_n]$

The Technique: Regularity

Set $S = \mathbb{R}[x_1, \dots, x_n]$

A graded S -module M has a **graded minimal free resolution**:

$$0 \rightarrow F_\delta \rightarrow F_{\delta-1} \rightarrow \cdots \rightarrow F_0 \rightarrow M \rightarrow 0, \quad \text{where } F_i \cong \bigoplus_j S(-a_{ij})$$

The Technique: Regularity

Set $S = \mathbb{R}[x_1, \dots, x_n]$

A graded S -module M has a **graded minimal free resolution**:

$$0 \rightarrow F_\delta \rightarrow F_{\delta-1} \rightarrow \cdots \rightarrow F_0 \rightarrow M \rightarrow 0, \quad \text{where } F_i \cong \bigoplus_j S(-a_{ij})$$

- **Projective dimension** $\text{pdim}(M) := \delta$
- **Castelnuovo-Mumford Regularity** $\text{reg}(M) := \max_{i,j} (a_{ij} - i)$

The Technique: Regularity

Set $S = \mathbb{R}[x_1, \dots, x_n]$

A graded S -module M has a **graded minimal free resolution**:

$$0 \rightarrow F_\delta \rightarrow F_{\delta-1} \rightarrow \cdots \rightarrow F_0 \rightarrow M \rightarrow 0, \quad \text{where } F_i \cong \bigoplus_j S(-a_{ij})$$

- **Projective dimension** $\text{pdim}(M) := \delta$
- **Castelnuovo-Mumford Regularity** $\text{reg}(M) := \max_{i,j} (a_{ij} - i)$
- Note: $M \cong \bigoplus_j S(-a_j) \implies \text{reg}(M) = \max\{a_j\}$

The Technique: Regularity

Set $S = \mathbb{R}[x_1, \dots, x_n]$

A graded S -module M has a **graded minimal free resolution**:

$$0 \rightarrow F_\delta \rightarrow F_{\delta-1} \rightarrow \cdots \rightarrow F_0 \rightarrow M \rightarrow 0, \quad \text{where } F_i \cong \bigoplus_j S(-a_{ij})$$

- **Projective dimension** $\text{pdim}(M) := \delta$
- **Castelnuovo-Mumford Regularity** $\text{reg}(M) := \max_{i,j} (a_{ij} - i)$
- Note: $M \cong \bigoplus_j S(-a_j) \implies \text{reg}(M) = \max\{a_j\}$

$\text{reg}(M)$ governs when $HF(M, d) = HP(M, d)$ [Eisenbud '05]:

$$HF(M, d) = HP(M, d) \text{ for } d \geq \text{reg}(M) + \text{pdim}(M) - n + 1$$

The Technique: Regularity

Set $S = \mathbb{R}[x_1, \dots, x_n]$

A graded S -module M has a **graded minimal free resolution**:

$$0 \rightarrow F_\delta \rightarrow F_{\delta-1} \rightarrow \cdots \rightarrow F_0 \rightarrow M \rightarrow 0, \quad \text{where } F_i \cong \bigoplus_j S(-a_{ij})$$

- **Projective dimension** $\text{pdim}(M) := \delta$
- **Castelnuovo-Mumford Regularity** $\text{reg}(M) := \max_{i,j} (a_{ij} - i)$
- Note: $M \cong \bigoplus_j S(-a_j) \implies \text{reg}(M) = \max\{a_j\}$

$\text{reg}(M)$ governs when $HF(M, d) = HP(M, d)$ [Eisenbud '05]:

$$HF(M, d) = HP(M, d) \text{ for } d \geq \text{reg}(M) + \text{pdim}(M) - n + 1$$

Results on previous slide follow from bounding $\text{reg}(C^r(\widehat{\Delta}))$.

Obtaining the Regularity Bound

Two key properties:

- 1 Regularity of any module in $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ can be bounded by regularity of other two.

Obtaining the Regularity Bound

Two key properties:

- 1 Regularity of any module in $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ can be bounded by regularity of other two.
- 2 If $A \subset B$ is a submodule and $\text{pdim}(B) < \text{codim}(B/A)$, then $\text{reg}(B) \leq \text{reg}(A)$.

Obtaining the Regularity Bound

Two key properties:

- 1 Regularity of any module in $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ can be bounded by regularity of other two.
 - 2 If $A \subset B$ is a submodule and $\text{pdim}(B) < \text{codim}(B/A)$, then $\text{reg}(B) \leq \text{reg}(A)$.
- Regularity bound obtained by finding an approximation $LS^{r,1}(\widehat{\Delta}) \subset C^r(\widehat{\Delta})$ satisfying property 2.

Obtaining the Regularity Bound

Two key properties:

- 1 Regularity of any module in $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ can be bounded by regularity of other two.
 - 2 If $A \subset B$ is a submodule and $\text{pdim}(B) < \text{codim}(B/A)$, then $\text{reg}(B) \leq \text{reg}(A)$.
- Regularity bound obtained by finding an approximation $LS^{r,1}(\widehat{\Delta}) \subset C^r(\widehat{\Delta})$ satisfying property 2.
 - $LS^{r,1}(\widehat{\Delta})$ is the subalgebra of $C^r(\widehat{\Delta})$ generated by splines supported on the union of two adjacent facets.

Obtaining the Regularity Bound

Two key properties:

- 1 Regularity of any module in $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ can be bounded by regularity of other two.
 - 2 If $A \subset B$ is a submodule and $\text{pdim}(B) < \text{codim}(B/A)$, then $\text{reg}(B) \leq \text{reg}(A)$.
- Regularity bound obtained by finding an approximation $LS^{r,1}(\widehat{\Delta}) \subset C^r(\widehat{\Delta})$ satisfying property 2.
 - $LS^{r,1}(\widehat{\Delta})$ is the subalgebra of $C^r(\widehat{\Delta})$ generated by splines supported on the union of two adjacent facets.
 - Property 1 used to break bounding $\text{reg}(LS^{r,1}(\widehat{\Delta}))$ down into a local problem by fitting into exact complexes.

Obtaining the Regularity Bound

Two key properties:

- 1 Regularity of any module in $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ can be bounded by regularity of other two.
 - 2 If $A \subset B$ is a submodule and $\text{pdim}(B) < \text{codim}(B/A)$, then $\text{reg}(B) \leq \text{reg}(A)$.
- Regularity bound obtained by finding an approximation $LS^{r,1}(\widehat{\Delta}) \subset C^r(\widehat{\Delta})$ satisfying property 2.
 - $LS^{r,1}(\widehat{\Delta})$ is the subalgebra of $C^r(\widehat{\Delta})$ generated by splines supported on the union of two adjacent facets.
 - Property 1 used to break bounding $\text{reg}(LS^{r,1}(\widehat{\Delta}))$ down into a local problem by fitting into exact complexes.
 - Local problem solved directly

Two other applications of algebraic techniques:

- Analogue of basis with local support for nonsimplicial Δ [D. '14]

Two other applications of algebraic techniques:

- Analogue of basis with local support for nonsimplicial Δ [D. '14]
- Bounds on $\dim C_d^r(\Delta)$, $\Delta \subset \mathbb{R}^2, \mathbb{R}^3$ simplicial [Mourrain-Villamizar '13, Mourrain-Villamizar '14] - latter involves problem of **fat points** in \mathbb{P}^2

Two other applications of algebraic techniques:

- Analogue of basis with local support for nonsimplicial Δ [D. '14]
- Bounds on $\dim C_d^r(\Delta)$, $\Delta \subset \mathbb{R}^2, \mathbb{R}^3$ simplicial
[Mourrain-Villamizar '13, Mourrain-Villamizar '14] - latter involves problem of **fat points** in \mathbb{P}^2

Main Problem:

Planar case: Lower existing regularity bounds!

Two other applications of algebraic techniques:

- Analogue of basis with local support for nonsimplicial Δ [D. '14]
- Bounds on $\dim C_d^r(\Delta)$, $\Delta \subset \mathbb{R}^2, \mathbb{R}^3$ simplicial [Mourrain-Villamizar '13, Mourrain-Villamizar '14] - latter involves problem of **fat points** in \mathbb{P}^2

Main Problem:

Planar case: Lower existing regularity bounds!

Planar simplicial case: Show $\dim C_d^r(\Delta) = HP(C^r(\widehat{\Delta}), d)$ for $d \geq 2r + 1$.

- Regularity techniques in [D. '14] give equality in simplicial case for $d \geq 3r + 2$ (one off from Alfeld-Schumaker result).
- [Schenck-Stiller '02] use vector bundle techniques on projective space to approach regularity of $C^r(\widehat{\Delta})$.

Thank You!

References I



P. Alfeld, L. Schumaker, *On the dimension of bivariate spline spaces of smoothness r and degree $d = 3r + 1$* , Numer. Math. **57** (1990) 651-661.



L. Billera, *Homology of Smooth Splines: Generic Triangulations and a Conjecture of Strang*, Trans. Amer. Math. Soc. **310**, 325-340 (1988).



L. Billera, *The Algebra of Continuous Piecewise Polynomials*, Adv. in Math. **76**, 170-183 (1989).



L. Billera, L. Rose, *A Dimension Series for Multivariate Splines*, Discrete Comput. Geom. **6**, 107-128 (1991).



L. Billera, L. Rose, *Modules of piecewise polynomials and their freeness*, Math. Z. **209** (1992), 485-497.



M. DiPasquale, *Shellability and Freeness of Continuous Splines*, J. Pure Appl. Algebra **216** (2012) 2519-2523.



M. DiPasquale, *Lattice-Supported Splines on Polytopal Complexes*, Adv. in Appl. Math. **55** (2014) 1-21.



M. DiPasquale, *Regularity of Mixed Spline Spaces*, submitted.



D. Eisenbud, *The Geometry of Syzygies*, Springer-Verlag, New York, 2005.



G. Farin, *Curves and Surfaces for Computer Aided Geometric Design*, 4th ed., Academic Press, Boston, 1997.



D. Hong, *Spaces of bivariate spline functions over triangulation*, Approx. Theory Appl. **7** (1991), 56-75.



A. Ibrahim, L. Schumaker, *Super spline spaces of smoothness r and degree $d \geq 3r + 2$* , Constr. Approx. **7** (1991), 401-423.



T. McDonald, H. Schenck, *Piecewise Polynomials on Polyhedral Complexes*, Adv. in Appl. Math. **42**, no. 1, 82-93 (2009).

References II



B. Mourrain, N. Villamizar, Homological techniques for the analysis of the dimension of triangular spline spaces, *J. Symbolic Comput.* 50 (2013), 564-577.



B. Mourrain, N. Villamizar, Bounds on the dimension of trivariate spline spaces: A homological approach, *Math. Comput. Sci.* 8 (2014), 157-174.



S. Payne, *Equivariant Chow Cohomology of Toric Varieties*, *Math. Res. Lett.* 13, 29-41 (2006).



H. Schenck, Homological methods in the theory of splines, Thesis, Cornell University (1997).



H. Schenck, M. Stillman, *Local Cohomology of Bivariate Splines*, *J. Pure Appl. Algebra* 117 & 118, 535-548 (1997).



H. Schenck, A spectral sequences for splines, *Adv. in Appl. Math.* 19, 183-199 (1997).



H. Schenck, P. Stiller, Cohomology vanishing and a problem in approximation theory, *Manuscripta Math.* 107 (2002), 43-58.



G. Strang, *Piecewise Polynomials and the Finite Element Method*, *Bull. Amer. Math. Soc.* 79 (1973), 1128-1137.



S. Yuzvinsky, Modules of splines on polyhedral complexes, *Math. Z.* 210 (1992), 245-254.