Math 4513 Solutions to Homework Set 1

1.1. Given that

$$\frac{d^n}{dx^n} \left(\ln |x| \right) = (-1)^{n-1} (n-1)! x^{-n+1}$$

(a) Use the Taylor Theorem with Integral Remainder to find the magnitude of the error term $R_{100}(1.99)$ when one approximates $\ln[1.99]$ using the first 101 terms of the Taylor expansion about 1 of $\ln |x|$.

• We have

$$\ln|x| = \sum_{n=0}^{N} \frac{1}{n!} \frac{d^n}{dx^n} (\ln|x|) \bigg|_{x=1} (x-1)^n + R_N(x)$$

When N = 100, the error term is $R_{100}(x)$ given by

$$R_{100}(x) = \frac{1}{100!} \int_{1}^{x} \frac{d^{101}}{dt^{101}} \left(\ln |t| \right) \bigg|_{t} (x-t)^{100} dt$$
$$= \frac{1}{100!} \int_{1}^{x} (-1)^{100} (100!) t^{-100} (x-t)^{100} dt$$

For x = 1.99 we have

$$R_{100}(1.99) = \int_{1}^{1.99} t^{-100} (1.99 - t)^{100} dt$$

= -.4357523477 × 10⁻³⁵

(b) Use the Taylor Theorem with Lagrange Remainder to obtain an upper bound on the error term $E_{100}(x)$ when x ranges from 1.985 to 1.995 for the Taylor expansion of $\ln |x|$ about 1.

 $\bullet~$ We have

$$\ln|x| = \sum_{n=0}^{N} \frac{1}{n!} \left. \frac{d^n}{dx^n} \left(\ln|x| \right) \right|_{x=1} (x-1)^n + E_N(x)$$

When N = 100 the error term $E_{100}(x)$ is given by

$$E_{100}(x) = \frac{1}{(101)!} \frac{d^{101}}{dt^{101}} (\ln|t|) \Big|_{t=\xi} (x-1)^{101}$$

= $\frac{1}{101!} ((-1)^{100} (100!) \xi^{-100}) (x-1)^{101}$
= $\frac{1}{101!} \xi^{-100} (x-1)^{101}$

for some $\xi \in [1, x]$. When x = 1.99 we have

$$E_{100}(1.99) = \frac{1}{101} (0.99)^{101} \xi^{-100}$$

Clearly, the maximal value of this error term on the interval [1, 1.99] occurs when $\xi = 1$. so

$$\max E_{100}(1.99) = \frac{1}{101} (0.99)^{101} = 3.5878 \times 10^{-3}$$

1.2. Show that if $b_n = \mathcal{O}(a_n)$ then $b_n / \ln |n| = \mathfrak{o}(a_n)$.

• Since $b_n = \mathcal{O}(a_n)$

 $b_n \leq C a_n$ for sufficiently large n

Let $c_n = b_n / \ln |n|$. To see if $c_n = \mathfrak{o}(a_n)$ we look at

$$\lim_{n \to \infty} \frac{|c_n|}{|a_n|}$$

Now

$$0 \leq \frac{|c_n|}{|a_n|} = \frac{|b_n|}{|a_n|\ln|n|} \leq \frac{C}{\ln|n|} \quad \text{for sufficient large } n$$

But now

$$\lim_{n \to \infty} \frac{C}{\ln |n|} = 0 \quad \Rightarrow \quad \lim_{n \to \infty} \frac{|c_n|}{|a_n|} = 0$$

so $c_n = \mathfrak{o}(a_n)$.

1.3. Show that if $b_n = \mathfrak{o}(a_n)$ then $b_n = \mathcal{O}(a_n)$, but that the converse is not true.

• Since $b_n = \mathfrak{o}(a_n)$

$$\lim_{n \to \infty} \frac{|b_n|}{|a_n|} = 0$$

By definition, this means that for any number C there exists a number N such that n > N implies $|b_n| / |a_n| < C$. Or

 $|b_n| < C |a_n|$ for sufficiently large n

Hence, $b_n = \mathfrak{o}(a_n)$ implies $b_n = \mathcal{O}(a_n)$.

To see that the converse is not true, take $b_n = a_n$. Then $b_n = \mathcal{O}(a_n)$ (since $b_n \leq (1)a_n$ for all n > 1). But

$$\lim_{n \to \infty} \frac{|b_n|}{|a_n|} = \lim_{n \to \infty} \frac{|a_n|}{|a_n|} = 1 \neq 0$$

so $b_n \neq \mathfrak{o}(a_n)$.

1.4. Show that if $b_n = \mathcal{O}(a_n)$ and $c_n = \mathcal{O}(a_n)$, then $b_n + c_n = \mathcal{O}(a_n)$.

• If $b_n = \mathcal{O}(a_n)$ and $c_n = \mathcal{O}(a_n)$, then there exist constants C_1, C_2 and numbers N_1, N_2 such that

But then

 $|b_n + c_n| \leq (C_1 + C_2) |a_n|$, for all *n* greater than $\max(N_1, N_2)$ Hence, by taking $C = C_1 + C_2$, $N = \max(N_1, N_1)$ we see that $b_n + c_n = \mathcal{O}(a_n)$.

1.5. Show that if $b_n = \mathfrak{o}(a_n)$ and $c_n = \mathfrak{o}(a_n)$, then $b_n + c_n = \mathfrak{o}(a_n)$.

• If $b_n = \mathfrak{o}(a_n)$ and $c_n = \mathfrak{o}(a_n)$,

$$0 = \lim \frac{|b_n|}{|a_n|} = \lim \frac{|c_n|}{|a_n|}$$

and so, using the triangle inequality and the fact that the limit of a sum of convergent sequences is equal to the sum of their limits we have

$$\lim_{n \to \infty} \frac{|b_n + c_n|}{|a_n|} \le \lim_{n \to \infty} \left(\frac{|b_n|}{|a_n|} + \frac{|c_n|}{|a_n|} \right) = \lim_{n \to \infty} \frac{|b_n|}{|a_n|} + \lim_{n \to \infty} \frac{|c_n|}{|a_n|} = 0 + 0 = 0$$

so $b_n + c_n = \mathfrak{o}(a_n)$.

1.6. Show that for any r > 0, $x^r = \mathcal{O}(e^x)$ as $x \to \infty$.

• We need to show taht we can always find a constant C and a number N such that

$$|x^r| \le C |e^x| \quad , \quad \text{for all } x > N$$

Without loss of generality we can assume x is positive and greater than 1. If we take the log both sides of this inequality (which maintains the inequality since $\ln |x|$ is monotonically increasing for x > 1) we get

$$|r\ln|x| \le \ln|C|x|$$

 \mathbf{or}

$$C \ge r \frac{\ln|x|}{x}$$

Applying l'Hospital's rule to the right hand side one can see that

$$\lim_{x \to \infty} \frac{\ln |x|}{x} = \lim_{x \to \infty} \frac{\frac{1}{x}}{1} = \lim_{x \to \infty} \frac{1}{x} = 0$$

So for large enough x we can guarantee that $C \ge r \frac{\ln |x|}{x}$ for any constant C > 0. Hence we can maintain

 $|x^r| \leq C |e^x|$, for all sufficiently large x

so $x^r = \mathcal{O}(e^x)$.

- 1.7. Show that for any r > 0, $\ln |x| = \mathcal{O}(x^r)$ as $x \to \infty$.
 - We need to show that

$$\ln|x| \le C |x^r|$$

for all suficiently large x. Consider

$$\lim_{x \to \infty} \frac{\ln |x|}{x^r}$$

applying l'Hospital's rule

$$\lim_{x \to \infty} \frac{\ln |x|}{x^r} = \lim_{x \to \infty} \frac{\frac{1}{x}}{rx^{r-1}} = \lim_{x \to \infty} \frac{1}{r}x^{-r} = 0 \quad \text{if } r > 0$$

Therefore, we can make the ratio

$$\frac{\ln|x|}{x^r} < C$$

for any positive number C by taking x sufficiently large. Hence,

$$\ln |x| < Cx^r$$
 for x sufficiently large \Rightarrow $\ln |x| = \mathcal{O}(x^r)$

1.8. If 1/10 is correctly rounded to the normalized binary number $(1.a_1a_2...a_{23})_2 \times 2^m$, what is the roundoff error? What is the relative roundoff error?

• As demonstrated in class

Keeping only 23 bits past the decimal point we obtain the machine number that's obtained by rounding down is

$$q_{-} = (1 + 2^{-1} + 2^{-4} + 2^{-5} + 2^{-8} + 2^{-9} + 2^{-12} + 2^{-13} + 2^{-16} + 2^{-17} + 2^{-20} + 2^{-21}) * 2^{-4}$$

= (0.099999999406)₁₀

The next machine number up will be

We have

$$\frac{1}{10} - q_{-} = 5.94 \times 10^{-9}$$
$$\frac{1}{10} - q_{+} = 1.5 \times 10^{-9}$$

And so q_+ is the closest machine number, and the roundoff error is 1.5×10^{-9} . The relative roundoff error is

$$\frac{\left|\frac{1}{10} - q_{+}\right|}{\frac{1}{10}} = 1.5 \times 10^{-8}$$

1.9. Give examples of real numbers for which

$$fl(x \odot y) \neq fl(fl(x) \odot fl(y))$$

Illustrate this for all four arithmetic operators $(+, -, \times, \div)$ using a hypothetical machine with 5 decimal (base t0) digits.

• Addition: Let x = 1.000014 and y = 1.000014. Then

$$x + y = 1.000028$$

 and

$$fl(x) = fl(y) = 1.00001$$

$$fl(x) + fl(y) = 2.00002$$

$$fl(fl(x) + fl(y)) = 2.00002$$

but

$$fl(x + y) = 1.00003 \neq 2.00002 = fl (fl(x) + fl(y))$$

• Subtraction:

$$fl(x - y) = -0.00001 \neq 0.00000 = fl(fl(x) - fl(y))$$

• Multiplication

$$x = 1.000004$$

$$y = 1.000004$$

$$fl(x) = 1.00000$$

$$fl(y) = 1.00000$$

$$fl(x) \times fl(y) = 1.00000$$

$$fl(fl(x) \times fl(y)) = 1.00000$$

$$x \times y = 1.0000800016$$

$$fl(x \times y) = 1.00008$$

$$fl(x \times y) = 1.00008 \neq 1.00000 = fl(fl(x) \times fl(y))$$

• Division

$$\begin{array}{rcrcrcr} x & = & 0.0000125 \\ y & = & 1.000000 \\ f(x) & = & 0.00001 \\ fl(y) & = & 1.00000 \\ fl(y)/fl(x) & = & 1.00000 \times 10^5 \\ fl(fl(y)/fl(x)) & = & 1.00000 \times 10^5 \\ y/x & = & 80,0000 \\ fl(y/x) & = & 8.0 \times 10^4 \end{array}$$

$$fl(y/x) = 8.0 \times 10^4 \neq 1.00000 \times 10^5 = fl(fl(y)/fl(x))$$