## LECTURE 3

# Methods of Proof, cont'd

Last time we discussed some of the basic techniques for proving propositions. We began with the notion of a **direct proof** and one of its implementations, the **forward-backward method**. Then we discussed an alternative to the direct proof, **proof by contradiction**.

I would like to begin today's lecture with another example of proof by contradiction.

EXAMPLE 3.1. Use the method of proof by contradiction to show that the equation

$$x^2 + 2mx + 2n = 0$$

has no odd roots if m and n are odd.

This proposition is equivalent to one of the form  $P \Rightarrow Q$  if we take P to be the statement "m and n are odd integers", and Q to be the statement "there exists no odd integer x such that  $x^2 + 2mx + 2n = 0$ ". Not-Q is then the statement that there exists an odd integer x such that  $x^2 + 2mx + 2n = 0$ .

To apply the method of proof by contradiction, we suppose P and not-Q are true and look for a contradiction with known facts.

So let x be an odd integer satisfying

$$x^2 + 2mx + 2n = 0$$

and suppose both m and n are odd. Then

$$x^2 = -2(mx+n) \quad .$$

The right hand side, being a multiple of 2, is clearly even. So  $x^2$  is even. On the other hand, x is supposed to be odd. In your last homework assignment you proved (hopefully) that if x is odd, then  $x^2$  is necessarily odd. We have thus arrived at a contradiction with a known fact. Since P and not-Q can not be simultaneously true, we have

"*P* is true" 
$$\Rightarrow$$
 "not-*Q* is false"

or

 $P \Rightarrow Q$ .

1. The Contrapositive Method

The **contrapositive method** is a variation of the proof by contradiction method in which one tries to work forward from the hypothesis "not-Q is true" to conclude that "not-P" is true. For if we can show that

$$\operatorname{not-}Q \Rightarrow \operatorname{not-}P$$

then we can conclude that

 $P \Rightarrow Q$  .

The justification for conclusion runs as follows:

(i) "P is true" 
$$\Rightarrow$$
 "not-P is false"

(*ii*) "not-P is false" and "not-Q 
$$\Rightarrow$$
 not-P is true"  $\Rightarrow$  "not-Q is false"

(*iii*) "not-Q is false" 
$$\Rightarrow$$
 "Q is true"

EXAMPLE 3.2. Let's use the contrapositive method to prove

**PROPOSITION 3.3.** If n is an integer and  $n^2$  is an odd integer, then n is odd.

# proof.

In this proposition the hypothesis P is "n is an integer and  $n^2$  is an odd integer" and the conclusion Q we're trying to reach is "n is odd". So

not-Q = n is not even not-P = n is not an integer or  $n^2$  is not an odd integer

In figuring out not-P we used the rule that not-(A and B) = not-A or not-B. However, this makes our new conclusion more complicated then necessary.

Let's instead first rephrase the statement we're trying to prove as

Suppose n is an integer. Then  $n^2$  is an odd integer implies n is odd

or diagramatically

 $n \in \mathbb{Z} \Rightarrow (n^2 \text{ is odd} \Rightarrow n \text{ is odd})$ 

But then the contrapositive of the propositional statement in parentheses is

 $n \text{ is not-odd} \Rightarrow n^2 \text{ is not-odd}$ 

or

 $n ext{ is even } \Rightarrow n^2 ext{ is even}$ 

Since the contrapositive is true (we've shown this in an earlier example), the original propositional statement is true.

Here is the entire proof. The statement

"If n is an integer and  $n^2$  is an odd integer, then n is odd"

is equivalent to

" $n \in \mathbb{Z} \implies (n^2 \text{ is odd} \implies n \text{ is odd})$ "

which is equivalent to

" $n \in \mathbb{Z} \implies (n \text{ is not-odd} \implies n^2 \text{ is not odd})$ "

which is equivalent to

" $n \in \mathbb{Z} \implies (n \text{ is even} \implies n^2 \text{ is even})$ "

which has already been demonstrated to be true. So the original proposition is true

## 2. Proof by Construction

Another method of proof is particularly useful for proving statements involving existential quantifiers (e.g., "there exists at least one ...").

This method works as follows:

In order to prove a statement of the form

"If such and such, then there exists an object such that so and so."

it suffices to construct (guess, produce, devise an algorithm to produce, etc.), using the hypothesis " *such* and *such*", the *object* in the conclusion and show that it satisfies the properties " *so and so*".

EXAMPLE 3.4. Prove that if a < b, then there exists a real number c such that

a < c < b .

Proof. Set

$$c = \frac{a+b}{2}$$

Then

$$c-a=\frac{a+b-2a}{2}=\frac{b-a}{2}>0$$

since b > a; and

$$b - c = \frac{2b - a - b}{2} = \frac{b - a}{2} > 0$$

for the same reason. Thus, we have constructed a number c with the desired properties. EXAMPLE 3.5.

**PROPOSITION 3.6.** If a, b, c, d, e and f are real numbers such that

$$(ad - bc) \neq 0 \quad ,$$

then the two equations

$$\begin{array}{rcl} ax + by & = & e \\ cx + dy & = & f \end{array}$$

can be solved for x and y.

proof. Our hypothesis is that a, b, c, d, e and f are real numbers such that

$$(ad - bc) \neq 0$$

and the conclusion we are trying to prove is that there exists real numbers x and y such that

$$\begin{array}{rcl} (3.1) & & ax+by &=& e \\ & & cx+dy &=& f \end{array} .$$

Our basic plan is to apply a little high school algebra to construct a solution of the systems of equations. Thus, we might solve the first equation for y and then substitute the resulting expression for y into the second equation and solve for x. But note at that first step we end up with

$$y = \frac{e - ax}{b}$$

which requires an additional hypothesis,  $b \neq 0$ , which is distinct from our original hypothesis  $ad - bc \neq 0$ . This being the case, to get an air-tight proof, we'll need to handle the two cases  $b \neq 0$  or b = 0 separately,

Case 1.  $b \neq 0$ . In this case, we can solve the first equation of y, to get

$$y = \frac{e - ax}{b}$$

Substituting this into the second equation we get

$$cx + d\left(\frac{e - ax}{b}\right) = f.$$

Multiplying this last equation by -b we get

$$-cbx - de + adx = -fb$$

or

$$(ad - bc) x = de - fb$$

Since by hypothesis,  $ad - bc \neq 0$ , we can divide both sides by ad - bc to get

$$x = \frac{de - fb}{ad - bc}$$

and then substituting this expression for x into our expression for y we get

$$y = \frac{af - ce}{ad - bc}$$

and so we arrive at a solution.

Case 2. If b = 0. In this situation, our equations specialize to

$$\begin{array}{rcl} ax & = & e \\ bx + dy & = & f \end{array}$$

Now, also in this case, it must be that neither a or d equals zero; otherwise we will violate our hypothesis that  $ad - bc \neq \dot{0}$ . We can therefore solve the first equation for x

$$x = \frac{e}{a} \quad ,$$

Substituting this expression for x into the second equation, we get

$$dy = f - \frac{e}{a}$$

Since  $d \neq 0$  in this case, we can obtain

$$y = \frac{1}{d} \left( f - \frac{e}{a} \right)$$

and we have a solution for this case.

Since Case 1 and Case 2 exhaust all possibilities (for b), our proof is complete